Rail

Solution

Way: Ad Hoc

Query complexity: 3(N-1)Time complexity: $O(N \log N)$

** This problem can be solved by the following steps :

First, we know that station 0 is C type, and its location. We can query all the other stations' distances from station 0, we call this: dis[0][i]

Second, we sort all the dis[0][i], and obviously, the station x with the shortest distance dis[0][x] must be the first D type location after station 0.

Third, we process each station one by one according to their shortest distance with station 0 (that is, the order obtained in second step). For each station processed, we determine its type and location immediately as follows:

- 3.1 First, we maintain the information (location,id) of the leftmost C type and the rightmost D type as the algorithm proceeds.
- 3.2 To process the current station k, we use two queries, query(k,leftmost C type) as dis[k][L], query(k, rightmost D type) as dis[k][R]. And we also have dis[0][k]. By some observations, we know that either dis[k][L] or dis[k][R] is achieved with a 'direct' route (without moving forth and back).

For example, we have only 4 cases to consider:

By careful analysis (some if/else conditions), we can get the answer. Sometimes, we may need extra information to resolve the four cases, where we can check with dis[0][k].

Take case (a) for example:

the location of k might be L + dis[k][L], Then use this location, we need to check if dis[k][R] is reasonable or not.

There might be some p(type C) between L and k, So the dis[k][R] is dis[p][R]+dis[p][k], we need to check whether p exists or not. If p doesn't exist, then case (a) might be wrong, then try cases (b),(c), and (d) until we find the answer.