

# High-Dimensional Geometry

Input file:            **standard input**  
Output file:           **standard output**  
Time limit:            4 seconds  
Memory limit:         1024 megabytes

When a geometric problem is extended to three dimensions or even higher dimensions, its difficulty can increase significantly. To address complex  $n$ -dimensional geometric structures, researchers have proposed a model in which the  $2^n$  key positions in space are abstracted as vertices of a graph, and specific geometric constraints between these positions are abstracted as edges in the graph.

You now have a model that has been abstracted into a graph, in which complex geometric features have been transformed into a simple undirected graph containing  $2^n$  vertices with no duplicate edges or self-loops. However, for certain reasons, you cannot directly observe the edges in the graph. You want to know the total number of edges,  $m$ .

You may use a “subset detection detector” which works as follows:

- You input a subset  $S$  containing  $k$  vertices into the detector. The detector requires that the size of each set examined must be **exactly** half the total number of vertices in the graph, i.e.,  $k = 2^{n-1}$ .
- The detector scans these vertices and their associated edges, and returns the total number of edges in the graph where **at least one endpoint** belongs to the set  $S$ .

Due to the high cost, the detector can be used at most  $2^{n+1}$  times. Given the limited number of queries, find the total number of edges  $m$  in the graph.

## Interaction Protocol

First, read an integer  $n$  from standard input ( $1 \leq n \leq 10$ ). This means the graph contains  $2^n$  vertices, numbered from 1 to  $2^n$ .

To make a query, output `? s` to standard output, where  $s$  is a binary string of length  $2^n$ . The  $i$ -th character should be:

- 1, if vertex  $i$  is included in  $S$ ;
- 0, otherwise.

You must ensure that the number of 1s in the string is exactly  $2^{n-1}$ . If you make an invalid query or exceed the query limit, the interactor will terminate your program.

After each query, you need to read an integer  $d$  from standard input, which denotes the number of edges that have at least one endpoint in the subset  $S$ .

To output the answer, output `! m` to standard output, where  $m$  is the total number of edges. After outputting the answer, you must terminate the program immediately.

The interactor is **non-adaptive**. This means that the graph is fixed before the interaction begins and does not change according to your queries.

**Note:** After each output, you must print a newline and flush the standard output buffer. Use the following methods to flush:

- For C or C++, use `fflush(stdout)` or `cout.flush()`
- For Java, use `System.out.flush()`
- For Python, use `stdout.flush()`

## Example

standard input	standard output
2	? 1010
1	? 0101
0	! 1