## Problem A. LaLa and Magic Circle (LiLi Version)

Input file:
Output file: standard output
Time limit: $\quad 10$ seconds
Memory limit: 1024 megabytes

## This is an output-only problem.

LaLa has a pile of magic circles in her laboratory.
A magic circle can be represented as a simple polygon drawn with special ink, and is usable if it is convex. i.e. all of its internal angles are equal or less than $\pi$.

LaLa plans to turn every magic circle into a usable one. However, they may lose all their magical power if done incorrectly. Thankfully, LaLa has the perfect magical tool for that.

The tool works as follows. When you toss in a magic circle, if the circle is usable, the tool reports that it is. Otherwise, it takes two distinct points $u$ and $v$ such that

- $u$ and $v$ lie on the boundary of the convex hull of the magic circle, and
- none of the points on the path from $u$ to $v$ through the boundary of the magic circle in counterclockwise order lie on the boundary of the convex hull of the magic circle, except for $u$ and $v$.

Then, it rotates the $u-v$ path by $\pi$ around the midpoint of $u$ and $v$. In other words, for each point $w$ on the $u-v$ path, $w$ becomes $u+v-w$ where the addition is done coordinate-wise over the two dimensional coordinate system over the paper the magic circle is drawn on. Note that the result of this modification is also a simple polygon.
Little did LaLa know, LaLa's sister, LiLi, overheard LaLa's plan. Knowing how lazy LaLa is, as a prank, LiLi will sneak in a magic circle that takes large amount of applications of the tool to make it usable. More specifically, LiLi will add a magic circle to the pile which is a union of equal or less than 1000 line segments and the tool can perform some sequence of modifications with between 120000 and 1000000 steps that turns it circle into a usable magic circle.
Write a program to help LiLi compute one such magic circle.

## Output

The output should be in the following format:

| $N$ |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
| $x_{0}$ | $y_{0}$ |  |  |  |
| $x_{1}$ | $y_{1}$ |  |  |  |
| $\vdots$ |  |  |  |  |
| $x_{N-1}$ | $y_{N-1}$ |  |  |  |
| $Q$ |  |  |  |  |
| $a_{0}$ | $b_{0}$ | $c_{0}$ | $d_{0}$ |  |
| $a_{1}$ | $b_{1}$ | $c_{1}$ | $d_{1}$ |  |
|  |  | $\vdots$ |  |  |
| $a_{Q-1}$ | $b_{Q-1}$ | $c_{Q-1}$ | $d_{Q-1}$ |  |

where the initial magic circle is the union of $N$ line segments connecting points $\left(x_{i}, y_{i}\right)$ and $\left(x_{(i+1 \bmod N)}, y_{(i+1 \bmod N)}\right)$ for all integers $0 \leq i<N$, and it has a sequence of modifications by the tool of length $Q$ such that $i$-th modification choose the counterclockwise path from point $\left(a_{i}, b_{i}\right)$ to $\left(c_{i}, d_{i}\right)$.

The output should satisfy the following constraints:

- All the numbers in the output are integers.
- $3 \leq N \leq 1000$
- $120000 \leq Q \leq 1000000$
- $0 \leq x_{i}, y_{i}, a_{j}, b_{j}, c_{j}, d_{j} \leq 1000000000$ for all integers $0 \leq i<N$ and $0 \leq j<Q$.
- $x_{i} \neq x_{j}$ or $y_{i} \neq y_{j}$ for all integers $0 \leq i<j<N$.
- $a_{i} \neq c_{i}$ or $b_{i} \neq d_{i}$ for all integers $0 \leq i<Q$.
- The sequence of points $\left(x_{0}, y_{0}\right),\left(x_{1}, y_{1}\right), \cdots,\left(x_{N-1}, y_{N-1}\right)$ defines a counterclockwise traversal of the boundary of a simple polygon.
- Choosing the counterclockwise path from $\left(a_{i}, b_{i}\right)$ to $\left(c_{i}, d_{i}\right)$ for the modification on the magic circle after the first $i$ modifications is valid for all integers $0 \leq i<Q$.
- The final magic circle is usable.

Note that it's allowed to have two consecutive segments meeting at the angle $\pi$. Also note that it is not required to minimize any number, your program just have to satisfy all the output constraints.

## Example

| standard input | standard output |
| :---: | :---: |
|  | $\begin{array}{lllll} \hline 8 & & & \\ 1 & 0 & & \\ 4 & 0 & & \\ 6 & 0 & & \\ 3 & 1 & & \\ 5 & 2 & & \\ 2 & 1 & & \\ 6 & 3 & & \\ 0 & 2 & & \\ 4 & & & \\ 6 & 0 & 6 & 3 \\ 10 & 2 & 6 & 3 \\ 10 & 2 & 9 & 4 \\ 9 & 4 & 0 & 2 \end{array}$ |

## Note

Please note that the sample output above does not satisfy the condition $120000 \leq Q \leq 1000000$, thus it will give Wrong Answer verdict upon submission. It is there only to present the output format.
The following illustrates the initial magic circle for the sample output.


The following illustrates the sequence of usage of the tool to make it usable. The dotted part of the boundary is the path modified by the tool, which becomes the red part after the modification.


## Problem B. LaLa and Magic Circle (LaLa Version)

Input file:
Output file:
Time limit:
Memory limit:
standard input standard output 2 seconds 1024 megabytes

LaLa has a pile of magic circles in her laboratory.
A magic circle can be represented as a simple polygon drawn with special ink, and is usable if and only if it is convex. i.e. all of its internal angles are equal or less than $\pi$.

LaLa plans to turn every magic circle into a usable one. However, it may lose all its magical power if done incorrectly. Thankfully, LaLa has the perfect magical tool for that.

The tool works as follows. When you toss in a magic circle, if it's usable, it reports that it is. Otherwise, it takes two distinct points $u$ and $v$ such that

- $u$ and $v$ lie on the boundary of the convex hull of the magic circle, and
- none of the points on the path from $u$ to $v$ through the boundary of the magic circle in counterclockwise order lie on the boundary of the convex hull of the magic circle, except for $u$ and $v$.

And then it rotates the $u-v$ path by $\pi$ around the midpoint of $u$ and $v$. In other words, for each point $w$ on the $u$-v path, $w$ becomes $u+v-w$ where the addition is done coordinate-wise over the two dimensional coordinate system over the paper the magic circle is drawn on. Note that the result of this modification is also a simple polygon.

LaLa got annoyed by how long it takes to convert them. In order to finish and take a nap ASAP, LaLa made the following observations.

1. A magic circle always turns into a usable one within a finite number of applications of the tool.
2. The set of points on the final magic circle is independent of the intermediate modifications. In other words, the shape and location of the final magic circle is a function of the initial magic circle.

Therefore, LaLa doesn't have to manually turn magic circles into usable ones with the tool. Instead, LaLa will compute the usable magic circle that can be made from the initial magic circle by a sequence of modifications by the tool and modify it in one go.
Write a program to help LaLa compute the final magic circle so that she can go take a nap.

## Input

The input is given in the following format:

| $N$ |  |
| :--- | :--- |
| $x_{0}$ | $y_{0}$ |
| $x_{1}$ | $y_{1}$ |
| $\vdots$ |  |
| $x_{N-1}$ | $y_{N-1}$ |

where the initial magic circle is the union of $N$ line segments connecting points $\left(x_{i}, y_{i}\right)$ and $\left(x_{(i+1 \bmod N)}, y_{(i+1 \bmod N)}\right)$ for all integers $0 \leq i<N$.
The input satisfies the following constraints:

- All numbers in the input are integers.
- $3 \leq N \leq 100000$
- $0 \leq x_{i} \leq 300000$ and $0 \leq y_{i} \leq 300000$ for all integers $0 \leq i<N$.
- $x_{i} \neq x_{j}$ or $y_{i} \neq y_{j}$ for all integers $0 \leq i<j<N$.
- The input defines a counterclockwise traversal of the boundary of a simple polygon. In particular, it does not intersect with itself.


## Output

The output should be in the following format:

| $M$ |  |
| :--- | :--- |
| $z_{0}$ | $w_{0}$ |
| $z_{1}$ | $w_{1}$ |
|  | $\vdots$ |
| $z_{M-1}$ | $w_{M-1}$ |

where the final usable magic circle is the union of $M$ line segments connecting points $\left(z_{i}, w_{i}\right)$ and $\left(z_{(i+1 \bmod M)}, w_{(i+1 \bmod M)}\right)$ for all integers $0 \leq i<M$,
The output should satisfy the following constraints:

- All the numbers in the output are integers.
- The point $\left(z_{0}, w_{0}\right)$ is lexicographically smaller than the point $\left(z_{i}, w_{i}\right)$ for all integers $1 \leq i<M$. i.e. $\left(z_{0}<z_{i}\right)$ or $\left(z_{0}=z_{i}\right.$ and $\left.w_{0}<w_{i}\right)$.
- Points $\left(z_{i}, w_{i}\right),\left(z_{(i+1 \bmod M)}, w_{(i+1 \bmod M)}\right)$, and $\left(z_{(i+2 \bmod M)}, w_{(i+2 \bmod M)}\right)$ are not collinear for all integers $0 \leq i<M$.
- The output defines a counterclockwise traversal of the boundary of a convex polygon.

It can be proved that the output satisfying the above constraints is unique.

## Example

$\left.\begin{array}{|ll|ll|}\hline & \text { standard input } & & \text { standard output } \\ \hline 8 & 0 & 6 & \\ 4 & 0 & 0 & 2 \\ 6 & 0 & 1 & 0 \\ 3 & 1 & 6 & 0 \\ 5 & 2 & 12 & 3 \\ 2 & 1 & 9 & 4 \\ 6 & 3 & 3 & 3\end{array}\right]$

## Note

The following illustrates the initial magic circle for the first sample.


The following illustrates the sequence of usage of the tool to make it usable. The dotted part of the boundary is the path modified by the tool, which becomes the red part after the modification.


## Problem C. LaLa and Lamp

| Input file: | standard input |
| :--- | :--- |
| Output file: | standard output |
| Time limit: | 3 seconds |
| Memory limit: | 1024 megabytes |

When LaLa laid down on her pet Leo's back to fall asleep, she noticed that the lamp is all messed up, which must have been the act of her sister LiLi.

The lamp can be modeled as a regular triangular grid where each cell contains a bulb which is either on or off.

LaLa wants to turn off the lamp (that is, set the state of all bulbs to off). LaLa can pick any of the three directions parallel to the side of a lamp, pick any row parallel to that direction, and then flip the state of all the bulbs in the row (on to off and off to on) with her magic. LaLa also could just walk over to the lamp and manually turn every bulb off, but she would prefer not to.

Write a program that determines whether LaLa can turn off the lamp with her magic.

## Input

The input is given in the following format:

| $N$ |
| :---: |
| $S_{0}$ |
| $S_{1}$ |
| $\vdots$ |
| $S_{N-1}$ |

where $N$ is the number of bulbs in a side of the lamp, and $S_{i}$ is the binary string of length $i+1$ representing the initial states of bulbs in the $i$-th row, where the $j$-th character of $S_{i}$ is ' 1 ' if and only if the $j$-th bulb is on.
The input satisfies the following constraint:

- $N$ is an integer.
- $2 \leq N \leq 2000$


## Output

If LaLa can turn off the lamp with magic, print a single string "Yes". Otherwise, print a single string "No". You may print each character in either case (lower or upper).

## Example

|  | standard input | standard output |
| :--- | :--- | :--- |
| 6 | Yes |  |
| 0 |  |  |
| 00 |  |  |
| 000 |  |  |
| 0110 |  |  |
| 00100 |  |  |

## Note

The following illustrates a sequence of magic LaLa should cast to turn off the lamp given in the sample. Empty circles denote the bulbs that are off, yellow circles denote the bulbs that are on, and red line is the choosen row for magic.

| $\begin{gathered} \text { stapo } \\ 00 \\ 000 \\ 0000 \\ 000000 \end{gathered}$ | $\begin{gathered} \operatorname{steq} 1 \\ 00 \\ 000 \\ 0000 \\ 00000 \\ 000000 \end{gathered}$ |  |
| :---: | :---: | :---: |
|  |  |  |

## Problem D. LaLa and Magic Stone

Input file:
Output file:
Time limit:
Memory limit:
standard input
standard output
1 second
1024 megabytes

One day, LaLa realized that she ran out of magic stone components for her magic tools. (Think of them as a battery in our world.) So LaLa rushed to a nearby store and bought a slab of magic stone.
LaLa wants to cut the slab into magic stone components. The slab consists of $N \times M$ cells. Unfortunately, some cells are incompatible with LaLa's magic tools.

The required magic stone component is a 7 -cell U-shaped piece.


As LaLa forgot to buy magic stone glue, LaLa can't merge smaller pieces to form the required shape.
Furthermore, since LaLa hates wasting magic stones, LaLa will be satisfied if and only if the slab is cut so that every single compatible cell belongs to a required shape.
Write a program that computes the number of ways to cut the slab so that LaLa is satisfied, modulo 998244353 . Two ways to cut pieces are different if and only if there exist two compatible cells such that they belong to the same piece in one and to different pieces in the other.

## Input

The input is given in the following format:

| $N$ | $M$ |
| :---: | :---: |
| $S_{0}$ |  |
| $S_{1}$ |  |
| $\vdots$ |  |
| $S_{N-1}$ |  |

where $N$ is the number of rows of the slab, $M$ is the number of columns, and $S_{i}$ is a binary string of length $M$ where $j$-th character is ' 1 ' if and only if the cell at the $i$-th row from the top and $j$-th column from the left is incompatible.

The input satisfies the following constraints:

- $N$ and $M$ are integers.
- $3 \leq N, M \leq 1000$


## Output

The output should be a integer equal to the number of ways to cut the slab so that LaLa is satisfied, modulo 998244353.

## Examples

|  | standard input | standard output |
| :--- | :--- | :--- |
| 44 | 2 |  |
| 0001 |  |  |
| 0000 |  |  |
| 1000 | 4 | 1 |
| 0001 |  |  |
| 0101 |  |  |
| 10000 |  |  |
| 1000 |  |  |

## Note

The following illustrates two possible ways to cut the slab in the first sample.


The following illustrates the only way to cut the slab in the second sample.


## Problem E. LaLa and Monster Hunting (Part 1)

Input file:
Output file:
Time limit:
Memory limit:
standard input
standard output
5 seconds
1024 megabytes

A dreadful monster has been witnessed in a forest near the city of magic Sharia, and a group of valorous adventurers will hunt it down in few days before it hurts anyone. However, LaLa knows that the real reason those adventurers are willing to take the risk is to obtain the rare magic stone that the monster is known to produce in its intestines. LaLa would like to obtain the magic stone before they do, as it is known to be quite beautiful.
LaLa will first locate the monster with her magic. LaLa has placed a bunch of magic tools within the forest, each of which has some power associated with it.
Consider the circles centered at each magic tool with radius equal to its power. LaLa's magic will successfully locate the monster if and only if the convex hull of the circles contains the location of the monster.

Write a program that determines whether LaLa will be able to locate the monster.

## Input

The input is given in the following format:

| $N$ |  |  |
| :--- | :---: | :---: |
| $x_{0}$ | $y_{0}$ | $r_{0}$ |
| $x_{1}$ | $y_{1}$ | $r_{1}$ |
|  | $\vdots$ |  |
| $x_{N-1}$ | $y_{N-1}$ | $r_{N-1}$ |

where $N$ is the number of magic tools placed in the forest, the $i$-th of which is located at $\left(x_{i}, y_{i}\right)$ and has power $r_{i}$. Here, assume that the forest is a two-dimensional plane where the monster is located at $(0,0)$.
The input satisfies the following constraints:

- All the numbers in the input are integers.
- $1 \leq N \leq 1000000$
- $-1000000 \leq x_{i}, y_{i} \leq 1000000$ for all integers $0 \leq i<N$
- $0 \leq r_{i} \leq 1000000$ for all integers $0 \leq i<N$
- The distance between point $(0,0)$ and the boundary of the convex hull of $N$ circles, $i$-th of which is centered at $\left(x_{i}, y_{i}\right)$ and has radius $r_{i}$, is at least 1 .


## Output

If LaLa's magic will successfully locate the monster, print a single string "Yes". Otherwise, print a single string "No". You may print each character in either case (lower or upper).

The 1st Universal Cup
Stage 14: Ranoa, Apr 29-30, 2023

## Examples

|  | standard input |  |  |
| :--- | :--- | :--- | :--- |
| 3 |  | Yes |  |
| -3 | 0 | 1 | standard output |
| 0 | 0 | 3 |  |
| 3 | 0 | 1 | Yes |
| 3 |  |  |  |
| 2 | 0 | 1 |  |
| 0 | 2 | 1 |  |
| -5 | -5 | 3 |  |
| 1 |  |  |  |
| 3 | 3 | 1 |  |

## Note

The following illustrates the configuration of the magic tools for the sample tests. The red curve denotes the boundary of the convex hull.
First Sample

## Problem F. LaLa and Monster Hunting (Part 2)

Input file:
Output file:
Time limit:
Memory limit:
standard input
standard output
6 seconds
1024 megabytes

A dreadful monster has been witnessed in a forest near the city of magic Sharia, and a group of valorous adventurers will hunt it down in few days before it hurt anyone. However, LaLa knows that the real reason those adventurers are willing to take the risk is to obtain the rare magic stone that the monster is known to produce in its intestines. LaLa would like to obtain the magic stone before they do, as it is known to be quite beautiful.
Currently, LaLa knows a rough estimate of the location of the monster. However, the monster excels at camouflage, so it's really hard to hunt it down when it's hiding in the network of branches.
For the sake of simplicity, we'll model the monster as a graph $G$ with 6 vertices described below:


The network of branches can be modeled as a simple graph $H$. A candidate is a subgraph of $H$ that is isomorphic to $G$. In other words, it is a graph obtained by deleting some edges from $H$, and then deleting some vertices that none of the remaining edges are incident to, whose vertices can be renumbered so that it coincides with $G$. LaLa will now have to examine all possible candidates to search and hunt the monster down.

Write a program that computes the number of candidates LaLa will have to examine, modulo 998244353.

## Input

The input describes the branch network $H$ and is given in the following format:

| $N$ | $M$ |
| :--- | :--- |
| $u_{0}$ | $v_{0}$ |
| $u_{1}$ | $v_{1}$ |
| $\vdots$ |  |
| $u_{M-1}$ | $v_{M-1}$ |

where $N$ is the number of joints, numbered from 0 to $N-1$ and $M$ is the number of branches, $i$-th of which connects the joints $u_{i}$ and $v_{i}$.
The input satisfies the following constraints:

- All the numbers in the input are integers.
- $2 \leq N \leq 100000$
- $0 \leq M \leq 100000$
- $0 \leq u_{i}<v_{i}<N$ for all integers $0 \leq i<M$
- $u_{i} \neq u_{j}$ or $v_{i} \neq v_{j}$ for all integers $0 \leq i<j<M$

Note that the network is not necessarily connected.

## Output

The output should be a single integer equal to the number of candidates, modulo 998244353.

## Examples

|  | standard input |  |
| :--- | :--- | :--- |
| 6 | 7 | standard output |
| 0 | 1 |  |
| 1 | 2 |  |
| 0 | 2 |  |
| 2 | 3 |  |
| 3 | 4 |  |
| 4 | 5 |  |
| 3 | 5 |  |
| 6 | 15 |  |
| 0 | 1 |  |
| 0 | 2 |  |
| 0 | 3 |  |
| 0 | 4 |  |
| 0 | 5 |  |
| 1 | 2 |  |
| 1 | 3 |  |
| 1 | 4 |  |
| 1 | 5 |  |
| 2 | 3 |  |
| 2 | 4 |  |
| 2 | 5 |  |
| 3 | 4 |  |
| 3 | 5 |  |
| 4 | 5 |  |

## Note

The followings illustrate the 4 candidates (the regular edges) of the branch network in the first sample test.


## Problem G. LaLa and Divination Magic

Input file:
Output file:
Time limit:
Memory limit
standard input
standard output
4 seconds
1024 megabytes

LaLa specializes in divination magic.
Let's say there are $M$ events $E_{0}, \cdots, E_{M-1}$ that LaLa's interested in forecasting. Each event is associated with one of two outcomes: catastrophe or salvation.
With a single use of LaLa's divination magic, LaLa obtains the knowledge of one of the following four forms:

1. Knowledge $(i, j, 1)$ : either $E_{i}$ is catastrophe or $E_{j}$ is catastrophe (possibly both).
2. Knowledge $(i, j, 2)$ : either $E_{i}$ is salvation or $E_{j}$ is catastrophe (possibly both).
3. Knowledge $(i, j, 3)$ : either $E_{i}$ is catastrophe or $E_{j}$ is salvation (possibly both).
4. Knowledge $(i, j, 4)$ : either $E_{i}$ is salvation or $E_{j}$ is salvation (possibly both).

LaLa cast her magic several times, possibly 0 , and wrote down all $M$-tuples of the outcomes of events that are consistent with her knowledge: this is called the result of the forecasting. And then, LaLa fell asleep.
When LaLa woke up, she found out that her pet, Leo, ruined all the predictions of her magic. Though LaLa was able to find the result of her forecasting, she is unsure if that data was ruined by Leo as well.
Write a program that determines whether there exists a set of predictions of LaLa's magic whose result of the forecasting matches the one LaLa has, and finds a possible set of predictions if there is one.

## Input

The input is given in the following format:

| $N$ | $M$ |
| :---: | :---: |
| $S_{0}$ |  |
| $S_{1}$ |  |
| $\quad \vdots$ |  |
| $S_{N-1}$ |  |

where $N$ is the number of outcomes in the result, $M$ of events, and $S_{i}$ is a binary string of length $M$ where $j$-th character is ' 1 ' if and only if the $i$-th result forecasts that $j$-th event will be in salvation.
The input satisfies the following constraints:

- $N$ and $M$ are integers.
- $1 \leq N, M \leq 2000$
- $S_{i} \neq S_{j}$ for all integers $0 \leq i<j<N$.


## Output

If there is no such prediction, the output should be a single integer -1 .
Otherwise, the output should be in the following format:

| $K$ |  |  |
| :--- | :---: | :---: |
| $I_{0}$ | $J_{0}$ | $t_{0}$ |
| $I_{1}$ | $J_{1}$ | $t_{1}$ |
|  | $\vdots$ |  |
| $I_{K-1}$ | $J_{K-1}$ | $t_{K-1}$ |

where $K$ is the size of a possible set $S$ of predictions, and, for each $0 \leq i<K, S$ contains the prediction Knowledge $\left(I_{i}, J_{i}, t_{i}\right)$.
The output should satisfy the following constraint:

- $0 \leq K \leq 2 \cdot M^{2}$

It can be proved that if there is such a set of predictions, there also is one satisfying the constraint.

## Examples

|  | standard input |  | standard output |  |
| :--- | :--- | :--- | :--- | :--- |
| 21 | 0 |  |  |  |
| 0 |  |  |  |  |
| 33 | 6 |  |  |  |
| 1011 | 0 | 2 | 3 |  |
| 111 | 0 | 1 | 4 |  |
|  | 0 | 2 | 4 |  |
| 1 | 2 | 3 |  |  |
|  | 1 | 2 | 4 |  |
|  | 2 | 2 | 4 |  |

## Problem H. LaLa and Harvesting

Input file:
Output file:
Time limit:
Memory limit:
standard input
standard output
4 seconds
1024 megabytes

Every other winter, LaLa helps her aunt, Aisha, harvest a crop native to the Biheiril Kingdom.
The crop can be modeled as a graph where an edge corresponds to a branch, a vertex to a joint, and a fruit with tastiness $T_{u}$ grows on each joint $u$.
The cultivation of the crop can be divided into three phases.

1. At the start of the first phase, a seed is planted on an open field. The seed eventually grows to form a cactus graph, a simple connected graph where every edge belongs to at most one cycle. This is the only phase where new joints grow.
2. Consider the DFS-tree of the cactus grown during the first phase. At the start of the second phase, Aisha encloses the crop with a ring-shaped framework by attaching the leaves (vertices of degree 1 , where only the edges in the DFS-tree are counted for degree, which possibly includes the root of the DFS-tree) to the ring. This forces the crop to grow additional branches joining adjacent joints on the framework, forming a cycle graph, and to start producing the fruits. Please read the input specification for detailed description.
3. At the start of the third phase, Aisha connects to the framework a magic tool that feeds the crop with the constant flow of extra nutrition. As a result, some new branches grow. The union of the branches grown during this phase form a very dense tree: every non-leaf vertex of the tree has degree $\geq 12$. The harvesting begins at the end of this phase.

LaLa's goal is to maximize the sum of tastiness of the fruits she has harvested. However, LaLa is not allowed to harvest fruits at two adjacent vertices, as it will stress and kill all the trunks directly connecting them.
Write a program that computes a set of fruits LaLa can harvest which has maximum possible sum of tastiness.

## Input

The input describes the state of a crop and is given in the following format:

| $N$ | $M$ |  |  |
| :--- | :--- | :--- | :--- |
| $T_{0}$ | $T_{1}$ | $\cdots$ | $T_{N-1}$ |
| $u_{0}$ | $v_{0}$ |  |  |
| $u_{1}$ | $v_{1}$ |  |  |
| $\quad \vdots$ |  |  |  |
| $u_{M-1}$ | $v_{M-1}$ |  |  |
| $K$ |  |  |  |
| $x_{0}$ | $y_{0}$ |  |  |
| $x_{1}$ | $y_{1}$ |  |  |
|  | $\vdots$ |  |  |
| $x_{K-1}$ | $y_{K-1}$ |  |  |

where $N$ is the number of joints, numbered from 0 to $N-1, M$ the number of branches grown during the first phase, and the $i$-th branch connects the joints $u_{i}$ and $v_{i}$ for all integers $0 \leq i<M$.
Consider the DFS-tree built over the cactus graph grown during the first phase where the DFS traversal starts at joint 0 and the order of the neighbors of each joints is given by the input order i.e. the traversal
prioritizes branches given sooner in the input. Note that the above description uniquely defines a DFStraversal and its corresponding DFS-tree. Let $c_{0}, \cdots, c_{l-1}$ be the subsequence of the DFS-order consisting of all joints which has degree 1 in the DFS-tree. Then the cycle graph grown during the second phase is defined by the $l$ edges $\left(c_{0}, c_{1}\right), \cdots,\left(c_{l-1}, c_{0}\right)$.

In addition, $K$ is the number of branches grown on the third phase, $i$-th of which connects the joints $x_{i}$ and $y_{i}$ for all integers $0 \leq i<K$.

The input satisfies the following constraints:

- All the numbers in the input are integers.
- $2 \leq N \leq 500, N-1 \leq M \leq 2 N, 1 \leq K \leq \min (N-1,100)$
- $1 \leq T_{u} \leq 200000$ for all integers $0 \leq u<N$
- $0 \leq u_{i}<v_{i}<N$ for all integers $0 \leq i<M$
- $u_{i} \neq u_{j}$ or $v_{i} \neq v_{j}$ for all integers $0 \leq i<j<M$
- $0 \leq x_{i}<y_{i}<N$ for all integers $0 \leq i<K$
- $x_{i} \neq x_{j}$ or $y_{i} \neq y_{j}$ for all integers $0 \leq i<j<K$
- The union of edges $\left(u_{i}, v_{i}\right)$ forms a cactus over the vertices $\{0,1, \cdots, N-1\}$.
- The union of edges $\left(x_{i}, y_{i}\right)$ forms a tree over the set of vertices $\left\{x_{0}, y_{0}, x_{1}, y_{1}, \cdots, x_{K-1}, y_{K-1}\right\}$ where each vertex with degree greater than 1 has degree at least 12 .


## Output

The output should be in the following format:

| $W$ | $L$ |  |  |
| :--- | :--- | :--- | :--- |
| $s_{0}$ | $s_{1}$ | $\cdots$ | $s_{L-1}$ |

where $W$ is the maximum sum of tastiness of a set $s=\left\{s_{0}, \cdots, s_{L-1}\right\}$ of fruits LaLa can harvest from the crop without harming any branches.
The output should satisfy the following constraints:

- All the numbers in the output are integers.
- $0 \leq L \leq N$
- $0 \leq s_{0}<s_{1}<\cdots<s_{L-1}<N$
- There are no edges directly connecting the vertices $s_{i}$ and $s_{j}$ for all integers $0 \leq i<j<L$.
- $\sum_{i=0}^{L-1} T_{s_{i}}=W$


## Example

| standard input | standard output |
| :---: | :---: |
| 67 | 22 |
| 111111 | 04 |
| 01 |  |
| 12 |  |
| 23 |  |
| 24 |  |
| 15 |  |
| 14 |  |
| 05 |  |
| 1 |  |
| 25 |  |

## Note

The following illustrates the crop described in the sample test.

- The regular black edges denote the branches grown during the first phase that is part of the DFStree.
- The dotted black edges denote the branches grown during the first phase that is NOT part of the DFS-tree
- The red edges denote the branches grown during the second phase.
- The blue edges denote the branches grown during the third phase.

Note that there can be multiple branches directly connecting the same pair of joints.


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## Problem I. LaLa and Spirit Summoning

Input file:
Output file:
Time limit:
Memory limit:
standard input
standard output
3 seconds
1024 megabytes

LaLa's younger sister LiLi is helping LaLa cast the spirit summoning magic.
While LaLa was asleep, LiLi had already built a prototype of the spirit to summon. The spirit consists of $N$ magic joints which allow any magic bars attached to them to freely move around them, and $M$ magic bars of various colors, each of which connects two magic joints and whose length can be adjusted to any non-negative real number before the summoning (but not after).

When it comes to spirit summoning, LaLa has a far higher standard than LiLi. Of course, LaLa was not satisfied with LiLi's work whatsoever. LaLa would like to fabulize the prototype by getting rid of some magic bars so that

1. the spirit is beautiful, which means there should not be two magic bars of the same color present, and
2. the spirit is as easy to control as possible, which means the degree of freedom of the spirit must be minimum over all beautiful spirits obtainable by eliminating some magic bars. Note that the minimum always exists as she can always eliminate all magic bars to create a beautiful spirit. See the note below for the exact definition of the degree of freedom.

Write a program that computes the degree of freedom of the spirit fabulized by LaLa.

## Input

The input describes the prototype spirit made by LiLi and is given in the following format:

| $N$ | $M$ |  |
| :--- | :---: | :---: |
| $u_{0}$ | $v_{0}$ | $c_{0}$ |
| $u_{1}$ | $v_{1}$ | $c_{1}$ |
|  | $\vdots$ |  |
| $u_{M-1}$ | $v_{M-1}$ | $c_{M-1}$ |

where $N$ is the number of magic joints, numbered from 0 to $N-1, M$ is the number of magic bars, and for each integer $0 \leq i<M$, the $i$-th magic bar has color $c_{i}$ and connects the magic joint $u_{i}$ and $v_{i}$.
The input satisfies the following constraints:

- All the numbers in the input are integers.
- $2 \leq N \leq 200$
- $0 \leq M \leq 1000$
- $0 \leq u_{i}<v_{i}<N$ and $0 \leq c_{i}<M$ for all integers $0 \leq i<M$

Note that there can be multiple magic bars connecting the same pair of magic joints.

## Output

The output should be a single integer equal to the degree of freedom of the spirit fabulized by LaLa.

## Examples

| standard input | standard output |
| :---: | :---: |
| $\begin{array}{lll} 3 & 3 & \\ 0 & 1 & 0 \\ 0 & 2 & 0 \\ 1 & 2 & 0 \end{array}$ | $5$ |
| $\begin{array}{lll} \hline 3 & 3 & \\ 0 & 1 & 0 \\ 0 & 2 & 1 \\ 1 & 2 & 2 \end{array}$ | $3$ |
| $\begin{array}{lll} \hline & 4 & \\ 0 & 1 & 0 \\ 1 & 2 & 1 \\ 2 & 3 & 2 \\ 0 & 3 & 3 \end{array}$ | $4$ |
| $\begin{array}{\|lll} \hline 5 & 4 & \\ 0 & 1 & 0 \\ 1 & 2 & 1 \\ 2 & 3 & 2 \\ 3 & 4 & 3 \end{array}$ | 6 |

## Note

Intuitively, the degree of freedom is the number of axis of motions preserving edge lengths of the spirit embedded on a plane.
More formally, let $E$ be an assignment of planar coordinates (we'll call this an embedding) to all magic joints of a spirit. Note that such an embedding can be identified with an element in $\mathbb{R}^{2 N}$ by concatenating all coordinates, where $N$ is the number of magic joints.
Let $C(E)$ be the set of embeddings continuously reachable from $E$ as an element of $\mathbb{R}^{2 N}$ while preserving edge lengths. i.e. for each element $E^{\prime}$ of $C(E)$ and each magic bars of the spirit connecting magic joints $u$ and $v$, the euclidean distance between $u$ and $v$ must be the same in $E$ and $E^{\prime}$.

The degree of freedom of $E$ is the minimum non-negative integer $k$ such that there exists a continuous bijective mapping $F: D \rightarrow C(E)$ where $D$ is a connected subset of $\mathbb{R}^{k}$.

The degree of freedom of a spirit is the maximum degree of freedom over all such embeddings $E$.
The following illustrate the spirit fabulized by LaLa along with one of the optimal embedding and the mapping $F$ for each sample tests in order.

1. $k=5, D=\mathbb{R}^{2} \times[0,2 \pi) \times \mathbb{R}^{2}$
$F:\left(x_{0}, x_{1}, x_{2}, x_{3}, x_{4}\right) \mapsto\left\langle\left(x_{0}, x_{1}\right),\left(x_{0}, x_{1}\right)+\left(\cos x_{2}, \sin x_{2}\right),\left(x_{3}, x_{4}\right)\right\rangle$
The following illustrates the 5 degrees of freedom associated with each variables.

2. $k=3, D=\mathbb{R}^{2} \times[0,2 \pi)$
$F:\left(x_{0}, x_{1}, x_{2}\right) \mapsto\left\langle\left(x_{0}, x_{1}\right),\left(x_{0}, x_{1}\right)+\left(\cos x_{2}, \sin x_{2}\right),\left(x_{0}, x_{1}\right)+\left(\cos \left(\frac{\pi}{3}+x_{2}\right), \sin \left(\frac{\pi}{3}+x_{2}\right)\right)\right\rangle$
The following illustrates the 3 degrees of freedom associated with each variables.

3. $k=4, D=\mathbb{R}^{2} \times(((0,2] \times[0,2 \pi)) \cup(\{0\} \times[0, \pi)))$
$F:\left(x_{0}, x_{1}, x_{2}, x_{3}\right) \mapsto\left\langle P_{0}, P_{0}+\frac{x_{2}}{2} P_{1}+\sqrt{1-\frac{x_{2}^{2}}{4}} P_{2}, P_{0}+x_{2} P_{1}, P_{0}+\frac{x_{2}}{2} P_{1}-\sqrt{1-\frac{x_{2}^{2}}{4}} P_{2}\right\rangle$
where $P_{0}=\left(x_{0}, x_{1}\right), P_{1}=\left(\cos x_{3}, \sin x_{3}\right)$ and $P_{2}=\left(\sin x_{3},-\cos x_{3}\right)$.
The following figure on the left illustrates the 4 degrees of freedom associated with each variables. Note that the motion associated with the variable $x_{2}$ is non-rigid. The one on the right illustrates the motion associated with $x_{2}$ in detail.

4. $k=6, D=\mathbb{R}^{2} \times[0,2 \pi)^{4}$
$F:\left(x_{0}, x_{1}, x_{2}, x_{3}, x_{4}, x_{5}\right) \mapsto\left\langle P_{0}, P_{1}, P_{2}, P_{3}, P_{4}\right\rangle$
where $P_{0}=\left(x_{0}, x_{1}\right)$ and $P_{i}=P_{i-1}+\left(\cos x_{i+1}, \sin x_{i+1}\right)$ for all integers $1 \leq i \leq 4$.
The following illustrates the 6 degrees of freedom associated with each variables.


## Problem J. LaLa and Magical Beast Summoning

Input file:
Output file:
Time limit:
Memory limit:
standard input standard output 5 seconds 1024 megabytes

LaLa is about to cast a magical beast summoning magic.
The first thing LaLa do is creating a summoning field, which has 3 constants associated with it: nullity $M$, elasticity $E$, and viscosity $V$. Such summoning field is denoted by $\mathcal{F}(M, E, V)$

A magical beast summoning magic is performed over a summoning cell within the summoning field, which is square-shaped and is associated with 3 variables: side length $L$, agility $A$, and intelligence $I$. Such summoning cell is denoted by $\mathcal{C}(L, A, I)$.
$\mathcal{C}(L, A, I)$ is in a null state if $L=0$. Otherwise, it is in a positive state.
The density of $\mathcal{C}(L, A, I)$ in positive state is defined as $(A \times I) / L^{2}$.
The problem of determining whether a magical beast summoning magic will succeed requires very heavy computation involving solving a system of 9999999999 -th order partial differential equations over 999999999999999 variables. Fortunately, LaLa already did all the math for you!

The magical beast summoning magic over $\mathcal{C}(L, A, I)$ within $\mathcal{F}(M, E, V)$ succeeds if and only if the function $\operatorname{valid}(M, E, V, L, A, I)$ defined by the pseudocode in the note section returns true. We'll call such summoning cell valid.

Sometimes, LaLa isn't satisfied with the set of summoning cells she has, and wants to generate new ones by combining them. The problem of determining the result of combination of two valid summoning cells $C_{0}=\mathcal{C}\left(L_{0}, A_{0}, I_{0}\right)$ and $C_{1}=\mathcal{C}\left(L_{1}, A_{1}, I_{1}\right)$ within $F=\mathcal{F}(M, E, V)$ requires another heavy computation, but thankfully, LaLa already did all the math for you again!

The result of combining two such cells $C_{0}$ and $C_{1}$ within $F$, denoted by $\operatorname{Combine}_{F}\left(C_{0}, C_{1}\right)$, is given by the function combine $\left(M, E, V, L_{0}, A_{0}, I_{0}, L_{1}, A_{1}, I_{1}\right)$ defined by the pseudocode in the note section, which returns a triple $L_{2}, A_{2}, I_{2}$ satisfying $\mathcal{C}\left(L_{2}, A_{2}, I_{2}\right)=$ Combine $_{F}\left(C_{0}, C_{1}\right)$. Here, it can be proved that Combine $_{F}\left(C_{0}, C_{1}\right)$ is also valid. Note that swapping the order of $C_{0}$ and $C_{1}$ affects the result.

The result of combining $K \geq 3$ cells $C_{0}, \cdots, C_{K-1}$ within $F$ is given recursively by

$$
\operatorname{Combine}_{F}\left(C_{0}, \cdots, C_{K-1}\right)=\operatorname{Combine}_{F}\left(\operatorname{Combine}_{F}\left(C_{0}, \cdots, C_{K-2}\right), C_{K-1}\right)
$$

For the sake of completeness, we define $\operatorname{Combine}_{F}(C)=C$.
LaLa is aware of a very special property about the combining operation that allows her to efficiently solve the range density query problem below. Can you figure it out?

You're given a summoning field $F=\mathcal{F}(M, E, V)$ and an array of $N$ valid summoning cells

$$
C_{0}=\mathcal{C}\left(L_{0}, A_{0}, I_{0}\right), \cdots, C_{N-1}=\mathcal{C}\left(L_{N-1}, A_{N-1}, I_{N-1}\right)
$$

within $F$. Write a program that processes the following two types of $Q$ queries:

- 1 i L A I
- Set $C_{i} \leftarrow \mathcal{C}(L, A, I)$.
- 21 r
- Let $R=$ Combine $_{F}\left(C_{l}, \cdots, C_{r-1}\right)$. If $R$ is in the null state, print a single integer -1 . Otherwise, print the density of $R$, modulo $M$. Here, an irreducible fraction $p / q$, where $p$ is a non-negative integer and $q$ is a positive integer not divisible by $M$, modulo $M$ is defined to be the unique integer $p \times q^{-1} \bmod M$ where $q^{-1}$ is the multiplicative inverse of $q$ modulo $M$. It can be proved that if $R$ is in the positive state, the denominator of the density of $R$ as an irreducible fraction is not divisible by $M$ within the constraints of this problem.


## Input

The input is given in the following format:

| $M$ | $E$ | $V$ |
| :--- | :---: | :---: |
| $N$ |  |  |
| $L_{0}$ | $A_{0}$ | $I_{0}$ |
|  | $\vdots$ |  |
| $L_{N-1}$ | $A_{N-1}$ | $I_{N-1}$ |
| $Q$ |  |  |
| $q_{0}$ |  |  |
| $\quad \vdots$ |  |  |
| $q_{Q-1}$ |  |  |

Here, $q_{i}$ denotes the $i$-th query, and is given in the format described in the statement.
The input satisfies the following constraints:

- All the numbers in the input are integers.
- $M$ is a prime such that $900000000 \leq M \leq 1000000000$
- $1 \leq E, V \leq 100$
- $1 \leq N, Q \leq 100000$
- $0 \leq L_{i}, A_{i}, I_{i}<M$ for all integers $0 \leq i<N$
- $\mathcal{C}\left(L_{i}, A_{i}, I_{i}\right)$ within $\mathcal{F}(M, E, V)$ is valid for all integers $0 \leq i<N$.
- For each query 1 i L A $\mathrm{I}, 0 \leq i<N, 0 \leq L, A, I<M$, and $\mathcal{C}(L, A, I)$ is valid within $\mathcal{F}(M, E, V)$.
- For each query 2 l r, $0 \leq l<r \leq N$


## Output

For each query of the second type, print its answer in a single line.

## Example

| standard input |  |  |  |
| :--- | :--- | :--- | :--- |
| 998244353 | 1 | 2 | standard output |
| 3 |  |  | 748683259 |
| 2 | 998244352 | 3 | 156877648 |
| 4 | 998244351 | 6 | 748683265 |
| 4 | 929561374 | 68682991 | 499122176 |
| 7 |  |  |  |
| 2 | 0 | 2 |  |
| 2 | 0 | 3 |  |
| 2 | 1 | 3 |  |
| 1 | 1 | 6 | 9 |
| 2 | 0 | 2 |  |
| 2 | 0 | 3 |  |
| 2 | 1 | 3 |  |

## Note

The following pseudocode defines the validity of summoning cells and the Combine operation.

## Both functions do not modify their arguments

| ```function \(\operatorname{valid}(M, E, V, L, A, I)\) if \(\min (L, A, I)<0 \quad\) or \(\quad M \leq \max (L, A, I)\) then return False end if /*Every operations and comparisons below are done mod \(M^{*} /\) if \(L=0\) and \((A+I \neq 0\) or \(A=I)\) then return False end if return \(A^{3}-A^{2} L+3 A^{2} I+E A L^{2}+L^{3} V+2 A L I+E L^{2} I+3 A I^{2}-L I^{2}+I^{3} \neq 0\) end function``` | ```function combine \(\left(M, E, V, L_{0}, A_{0}, I_{0}, L_{1}, A_{1}, I_{1}\right)\) Ensure \(\operatorname{VALID}\left(M, E, V, L_{0}, A_{0}, I_{0}\right)\) Ensure VALID \(\left(M, E, V, L_{1}, A_{1}, I_{1}\right)\) /*Every operations and comparisons below are done \(\bmod M^{*} /\) if \(L_{1}=0\) then return \(L_{0}, A_{0}, I_{0}\) end if if \(L_{0}=0\) then return \(L_{1}, I_{1}, A_{1}\) end if \(B_{0} \leftarrow\left(A_{0}+I_{0}\right) \cdot L_{1}, B_{1} \leftarrow\left(I_{1}+A_{1}\right) \cdot L_{0}\) \(C_{0} \leftarrow\left(A_{0}-I_{0}\right) \cdot L_{1}, C_{1} \leftarrow\left(I_{1}-A_{1}\right) \cdot L_{0}\) if \(B 0=B 1\) then if \(C_{0}+C_{1}=0\) then return \(0,3,-3\) end if Sum \(\leftarrow A_{0}+I_{0}\), Dif \(\leftarrow A_{0}-I_{0}\) \(B \leftarrow 3 \cdot\) Sum \(^{2}+E \cdot L_{0}^{2}\) \(C \leftarrow 2 \cdot\) Dif \(\cdot L_{0}\) \(D \leftarrow 2 \cdot C \cdot S u m \cdot\) Dif \(E \leftarrow B^{2}-2 \cdot D\) \(X \leftarrow C \cdot E, Y \leftarrow B \cdot(D-E)-2 \cdot C^{2} \cdot D i f^{2}\) return \(2 \cdot C^{3}, X+Y, X-Y\) else \(B \leftarrow B_{0}-B_{1}, C \leftarrow C_{0}-C_{1}, D=L_{0} \cdot L_{1}\) \(E \leftarrow C^{2} \cdot D-B^{2} \cdot\left(B_{0}+B_{1}\right)\) \(X \leftarrow B \cdot E, Y \leftarrow C \cdot\left(B_{0} \cdot B^{2}-E\right)-C_{0} \cdot B^{3}\) return \(2 \cdot B^{3} \cdot D, X+Y, X-Y\) end if end function``` |
| :---: | :---: |

The author has attached a C++ implementation which will get "Time Limit Exceeded" verdict upon submission, but it will always print the correct answer within finite time. You may reuse some part of the implementation on your submission. (If you're reading the printed version, you may find the attachment at the bottom of the statement on the eolymp site.)

