C: Cangaroo

Problem Author: Abe Wits

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Statistics: 11 submissions, 3 accepted, 8 unknown

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■ Number of can placements is $F_{m+1}$, the $(m+1)$ th Fibonacci number. Time complexity: $\mathcal{O}\left(n \cdot F_{m+1}^{2}\right)=\mathcal{O}\left(n \cdot 3.3^{m}\right)$ when using bitmasks.
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