## Problem J. Closest Pair Algorithm

| Input file: | standard input |
| :--- | :--- |
| Output file: | standard output |
| Time limit: | 10 seconds |
| Memory limit: | 512 mebibytes |

There is a classic problem about finding two closest points amongst a set of points on a plane. There is also a classic randomized algorithm to solve it, which goes like this:

```
rotate the plane around the origin by a random angle phi
let p be the array of points
sort p by x coordinate in increasing order
ans = INF;
for (int i = 0; i < n; i++) {
    for (int j = i - 1; j >= 0; j--) {
        if (p[i].x - p[j].x >= ans) break;
        ans = min(ans, dist(p[i], p[j]));
    }
}
```

Here "INF" is some number greater than all distances. The function "dist" returns Euclidean distance between the given points. Note that all $x$ coordinates are distinct after rotation with probability 1.
I know that the algorithm works well in practice, but how well exactly? I ask you to compute the expected number of calls of the function "dist", assuming that the angle "phi" is chosen uniformly at random from the range $[0 ; 2 \cdot \pi)$.

## Input

The first line contains a single integer $n(2 \leq n \leq 250)$ - the number of points.
The next $n$ lines contains coordinates of points, one per line.
All the coordinates are not greater than $10^{6}$ by absolute value.
It is guaranteed that all points are distinct. It is guaranteed that no three points lie on the same line.

## Output

Print one number - the expected number of calls of the function "dist".
Your answer is considered correct if its absolute or relative error does not exceed $10^{-6}$.
Formally, let your answer be $a$, and the jury's answer be $b$. Your answer is accepted if and only if $\frac{|a-b|}{\max (1,|b|)} \leq 10^{-6}$.

## Examples

|  | standard input | standard output |
| :--- | :--- | :--- |
| 4 |  | 5.0000000000000 |
| 0 | 0 |  |
| 1 | 1 |  |
| 1 | 1 | 2.5000000000000 |
| 3 |  |  |
| 0 | 0 |  |
| 0 | 1 | 1 |
| 1 | 0 |  |

