

## Problem B. Rectangle Tree

Input file: *standard input*  
Output file: *standard output*  
Time limit: 6 seconds  
Memory limit: 512 mebibytes

Mr. Peanutbutter has recently discovered a nice  $n \times n$  field covered with various crops. Diane told Mr. Peanutbutter that for generations this particular field is planted with crops using tree-like rectangle method. While Mr. Peanutbutter got distracted by a bird, Diane continued.

A *combinatorial rectangle* in the field is a subset of squares of the field of the form  $A \times B$  where  $A$  and  $B$  are subsets of the set  $\{0, \dots, n-1\}$ .

A *rectangle tree* is a rooted binary tree with  $k$  vertices with the following properties. Each vertex  $v$  of the tree is labeled with a combinatorial rectangle  $r(v) \subseteq \{0, \dots, n-1\} \times \{0, \dots, n-1\}$ . If  $s$  is an inner node of the tree, and  $c_1$  and  $c_2$  are its direct descendants, then their combinatorial rectangles form a partition of  $r(s)$ : formally,  $r(s) = r(c_1) \cup r(c_2)$  and  $r(c_1) \cap r(c_2) = \emptyset$ . A node cannot have only one direct descendant.

Let  $\text{Crop}(x, y)$  be the crop that grows on the square  $(x, y) \in \{0, \dots, n-1\} \times \{0, \dots, n-1\}$  of the field. A rectangle tree  $T$  with the root `Root` computes the crop types of the field if  $r(\text{Root}) = \{0, \dots, n-1\} \times \{0, \dots, n-1\}$  and for each leaf  $\ell$ , the combinatorial rectangle  $r(\ell)$  has exactly one type of crop growing on it: that is, for any two  $(x, y), (x', y') \in r(\ell)$ , we have  $\text{Crop}(x, y) = \text{Crop}(x', y')$ .

The *depth* of tree  $T$  is the largest distance between the root of  $T$  and a leaf of  $T$ . Here, distance stands for the number of edges in the shortest path between the vertices.

The *size* of tree  $T$  is the number of vertices in it.

You are given a rectangle tree  $T$  computing crop types `Crop`. Let the size of  $T$  be  $S$ . Construct another rectangle tree  $T'$  computing `Crop` such that its depth is at most  $3 \log_2 S$  and its size is at most  $5S$ .

### Input

The first line contains a single integer  $n$ , the size of the field ( $1 \leq n \leq 1000$ ). Each of the next  $n$  lines contain  $n$  integers describing the types of the crops. The  $j$ -th integer in the  $i$ -th row is the type of crop in the square  $(i, j)$ . All types are positive integers not exceeding  $10^7$ .

The next line contains a single integer  $S$ , the size of the rectangle tree ( $1 \leq S \leq 10\,000$ ,  $S \cdot n \leq 10^6$ ). The  $i$ -th of the next  $S$  lines contains the description of the  $i$ -th vertex of the tree. It contains several space-separated integers:  $p, m_1, m_2, a_1, \dots, a_{m_1}, b_1, \dots, b_{m_2}$ . Here,  $p \in \{0, 1, \dots, S-1\}$  is the number of the parent of vertex  $i$  (if  $i$  is the root, then  $p = i$ ), and the combinatorial rectangle corresponding to this vertex is  $r(i) = \{a_1, \dots, a_{m_1}\} \times \{b_1, \dots, b_{m_2}\}$ .

It is guaranteed that, if  $\ell$  is a leaf of the tree, then all types of crops in  $r(\ell)$  are the same. Additionally, for each inner vertex  $v$  with direct descendants  $c_1$  and  $c_2$ , the rectangles  $r(c_1)$  and  $r(c_2)$  form a partition of  $r(v)$ :  $r(v) = r(c_1) \cup r(c_2)$  and  $r(c_1) \cap r(c_2) = \emptyset$ . Finally, if  $i$  is the root, then  $r(i) = \{0, \dots, n-1\} \times \{0, \dots, n-1\}$ .

### Output

Print a tree of depth at most  $3 \log_2 S$  and of size at most  $5S$  such that it is also a rectangle tree that computes the given crops. The tree should be printed in the same format as the one given in the input.

## Example

standard input	standard output
3	7
1 1 2	0 3 3 0 1 2 0 1 2
1 1 2	0 1 3 2 0 1 2
2 2 2	1 1 2 2 0 1
5	1 1 1 2 2
0 3 3 0 1 2 0 1 2	0 2 3 0 1 0 1 2
0 3 2 0 1 2 0 1	4 2 1 0 1 2
0 3 1 0 1 2 2	4 2 2 0 1 0 1
1 2 2 0 1 0 1	
1 1 2 2 0 1	