## Problem K. NaN in a Heap

## Prerequisite: NaN

NaN (Not a Number) is a special floating-point value introduced by the IEEE 754 floating-point standard in 1985. The standard speficies that, when NaN is compared with a floating-point value x ( x can be positive, zero, negative, or even NaN itself), the following results should be returned.

| Comparison | $\mathrm{NaN} \geq \mathrm{x}$ | $\mathrm{NaN} \leq \mathrm{x}$ | $\mathrm{NaN}>\mathrm{x}$ | $\mathrm{NaN}<\mathrm{x}$ | NaN $=\mathrm{x}$ | NaN $\neq \mathrm{x}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Result | False | False | False | False | False | True |

## Prerequisite: Heap

A heap is a data structure which can be represented by a sequence with special properties. The following algorithm demonstrates how to insert $n$ floating-point values $a_{1}, a_{2}, \cdots, a_{n}$ into a min-heap $H$ in order, where $H$ is a sequence and is initially empty.

In the following algorithm, let $h_{i}$ be the $i$-th element in sequence $H$, and let $j / 2$ be the maximum integer $x$ satisfying $2 x \leq j$.

```
Algorithm 1 Heapify
    function \(\operatorname{Heapify}(A)\)
        Let \(H\) be an empty sequence.
        for \(i \leftarrow 1\) to \(n\) do \(\quad \triangleright n\) is the number of elements to be inserted into heap.
            Append \(a_{i}\) to the back of \(H\).
            \(j:=i\)
            while \(j>1\) do
                if \(h_{j}<h_{j / 2}\) then \(\quad \triangleright\) Recall that if \(h_{j}\) or \(h_{j / 2}\) is NaN, this expression will be false.
                    Swap \(h_{j}\) and \(h_{j / 2}\).
                    \(j:=j / 2\)
            else
                break
            end if
            end while
        end for
        return \(H\)
    end function
```


## Problem

Given an integer $n$, consider permutations of these $n$ elements: all integers from 1 to ( $n-1$ ) (both inclusive), as well as a NaN value. We say a permutation $P$ of these $n$ elements is a "heap sequence", if there exists a permutation $Q$ also of these $n$ elements satisfying $P=\operatorname{HEAPIFY}(Q)$.
We now randomly pick a permutation of these $n$ elements with equal probability (that is, the probability of a specific permutation to be picked is $\frac{1}{n!}$ ), calculate the probability that the picked permutation is a heap sequence.

## Input

There are multiple test cases. The first line of the input contains an integer $T\left(1 \leq T \leq 10^{3}\right)$ indicating the number of test cases. For each test case:
The first and only line contains an integer $n\left(1 \leq n \leq 10^{9}\right)$.

## Output

For each test case output one line containing the answer.
It can be proven that the answer is a rational number $\frac{p}{q}$. To avoid issues related to precisions, please output the integer $\left(p q^{-1} \bmod M\right)$ as the answer, where $M=10^{9}+7$ and $q^{-1}$ is the integer satisfying $q q^{-1} \equiv 1(\bmod M)$.

## Example

|  | standard input | standard output |
| :--- | :--- | :--- |
| 5 | 1 |  |
| 1 | 666666672 |  |
| 3 | 55555556 |  |
| 7 | 596445110 |  |
| 10 | 3197361 |  |
| 20221218 |  |  |

## Note

For the second sample test case, there are 4 heap sequences.

- $\{\operatorname{NaN}, 1,2\}=\operatorname{HEAPIFY}(\{\operatorname{NaN}, 1,2\})$.
- $\{\operatorname{NaN}, 2,1\}=\operatorname{HEAPIFY}(\{\mathrm{NaN}, 2,1\})$.
- $\{1, \operatorname{NaN}, 2\}=\operatorname{HEAPIFY}(\{1, \operatorname{NaN}, 2\})=\operatorname{HEAPIFY}(\{2, \operatorname{NaN}, 1\})$.
- $\{1,2, \operatorname{NaN}\}=\operatorname{HEAPIFY}(\{1,2, \operatorname{NaN}\})=\operatorname{HEAPIFY}(\{2,1, \operatorname{NaN}\})$.

So the answer is $\frac{4}{3!}=\frac{2}{3}$ in rational number. As $3 \times 333333336 \equiv 1(\bmod M)$, we should output $2 \times 333333336 \bmod M=666666672$.

