

Problem B. Disjoint Set Union

Input file: **standard input**
Output file: **standard output**
Time limit: 4 seconds
Memory limit: 1024 megabytes

Recently, Little Cyan Fish has been learning about the Disjoint Set Union (DSU) data structure. It is a powerful data structure that allows you to add edges to a graph and test whether two vertices of the graph are connected.

The DSU maintains a rooted forest structure consisting of n vertices. Each vertex x ($1 \leq x \leq n$) has a unique parent $f[x]$. If $x = f[x]$, then x is the root of its subtree. Initially, each vertex forms a single rooted tree. That is, $f[x] = x$ for all $1 \leq x \leq n$.

The basic interface of DSU consists of these two operations:

- **find** x : returns the root of the tree where x is located.
- **unite** x y : let $x' \leftarrow \mathbf{find}(x)$ and $y' \leftarrow \mathbf{find}(y)$. If $x' = y'$, do nothing. Otherwise, modify the parent of x' to y' .

To speed up the **unite** operation, Little Cyan Fish uses an optimization called *Path Compression*:

- If we call **find**(x) for some vertex x , we set the parent of each vertex from x to the root directly to the root.

The following pseudocode describes the details of the DSU.

Algorithm 1 An implementation of DSU with Path Compression

```

1: procedure FIND( $f, x$ )
2:   if  $x = f[x]$  then
3:     return  $x$ 
4:   end if
5:    $f[x] \leftarrow \mathbf{find}(f, f[x])$ 
6:   return  $f[x]$ 
7: end procedure
8: procedure UNITE( $f, x, y$ )
9:    $x \leftarrow \mathbf{FIND}(f, x)$ 
10:   $y \leftarrow \mathbf{FIND}(f, y)$ 
11:  if  $x \neq y$  then
12:     $f[x] \leftarrow y$ 
13:  end if
14: end procedure
```

Little Cyan Fish loves the DSU very much, so he would like to play with it. He got an array f of length n , where $f[i] = i$ in the beginning. Then, Little Cyan Fish did the following operations many times (possibly zero):

- Choose an integer $1 \leq x \leq n$, apply **FIND**(x).
- Choose two integers $1 \leq x \leq n$ and $1 \leq y \leq n$, apply **UNITE**(x, y).

He will give you the array f after all his operations. However, you would like to transform the array f into another given array g by using the DSU operations described above. You are wondering if it is possible to apply any additional operations so that $f[i] = g[i]$ for all $1 \leq i \leq n$.

Input

There are multiple test cases. The first line contains one integer T ($1 \leq T \leq 10^5$), representing the number of test cases.

For each test case, the first line contains one positive integer n ($3 \leq n \leq 1\,000$).

The next line contains n integers f_1, f_2, \dots, f_n denoting the array f after Little Cyan Fish's operations. It is guaranteed that the array f can be generated by using the operation above.

The following line contains n integers g_1, g_2, \dots, g_n denoting the array g that you would like to transform the array f into.

It is guaranteed that the sum of n^2 over all test cases does not exceed 5×10^6 .

Output

For each test case, if it is impossible to transform the array f into the array g , print a single line NO.

Otherwise, the first line of the output should contain a single word YES.

The next line of the output should contain an integer m ($0 \leq m \leq 2 \cdot n^2$), indicating the number of operations you used.

The next m lines describe the operations you used. Each operation is described in the following format:

- 1 x : Call FIND(x).
- 2 x y : Call UNITE(x, y).

If there are multiple solutions, you may print any of them. It can be proved that if any solution exists, then there's a plan consisting of no more than $2 \cdot n^2$ operations.

Example

standard input	standard output
5	YES
3	1
1 2 3	2 1 2
2 2 3	YES
4	4
1 2 3 3	2 3 2
1 1 1 2	1 4
5	2 2 1
1 2 3 4 5	1 3
2 3 4 5 5	YES
5	4
1 1 1 1 1	2 1 2
1 2 3 4 5	2 1 3
6	2 2 4
1 2 2 4 5 6	2 3 5
1 1 5 1 4 2	NO
	YES
	7
	2 6 2
	2 2 5
	1 3
	2 2 4
	1 2
	2 2 1
	1 2